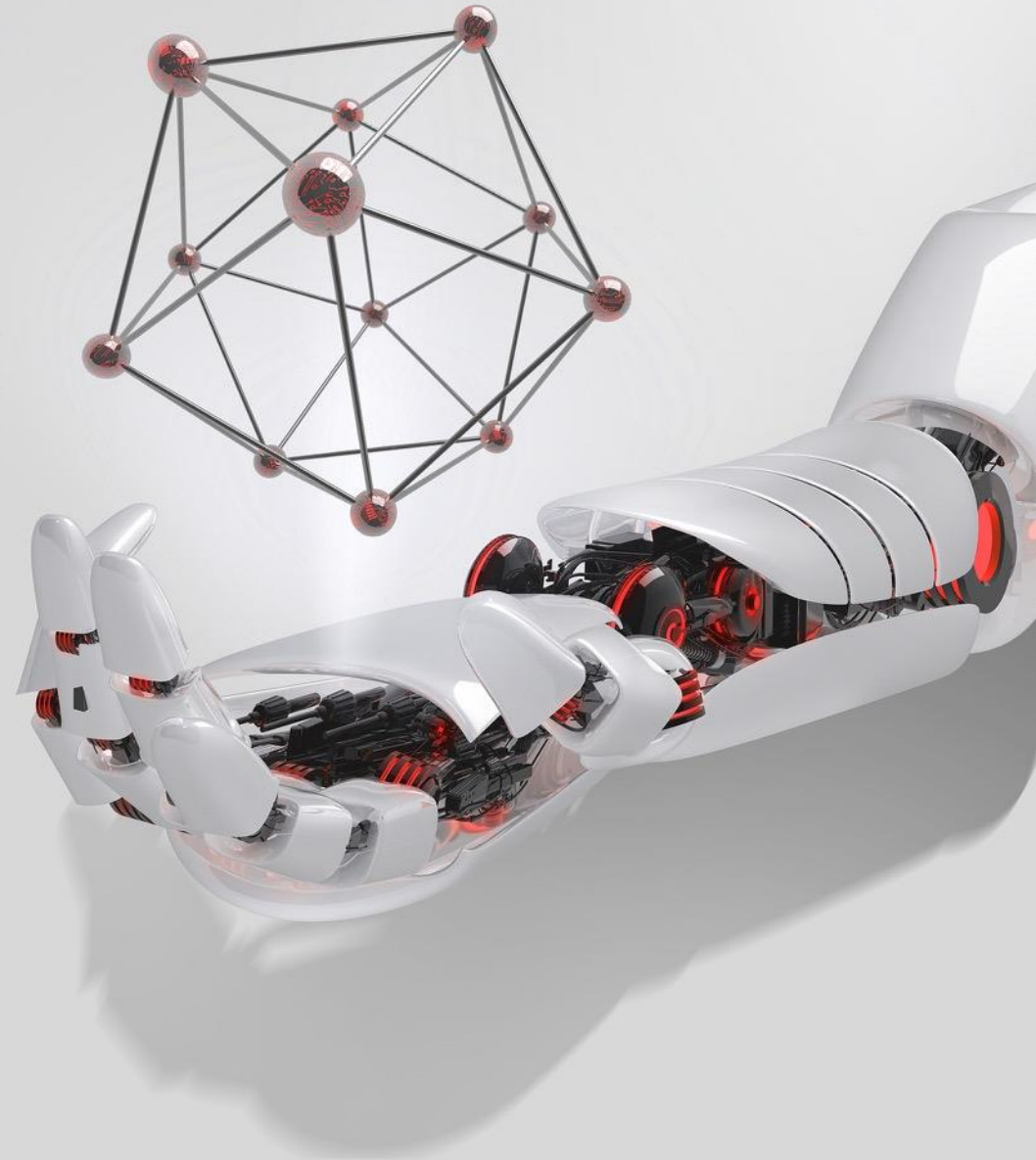


Decentralized Distributed Consensus at Expenses of Diffusion

David Guzman, Dirk Trossen, Jörg Ott

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Raitenhaslach

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Framing our Discussion Today

- **Distributed consensus systems (DCSes)** foresee participating peers to achieve consensus over a shared state in a continuous computation
 - **Majority rule** [VonNewman1956]
- We discuss a randomized consensus **approach**: why and how
- We present **empirical** insights
- We then discuss its implications for these **DCSes**: are they really **D**?
- We ask for further solutions to current approaches

The Consensus

- Distributed consensus arises in **replicated** state machines
 - State machines in a collection of **peers** compute **identical** copies of the **distributed** state and can **continue** operating even if some of the peers are untrusted (permissionless) or are down [Ongaro2014-USENIX]
- DCSes, e.g., DLTs:
 - Peers **replicate** information, i.e., transactions or blocks
 - **The system picks a random replica (peer). Decentralized**
 - The random peer **decides** on the distributed consensus.
 - This peer **replicates** (~ broadcasts) the consensus
- Information **replication** at the Internet scale is **key**
 - Current realization: randomized iterative diffusion

The Fault Tolerant Consensus

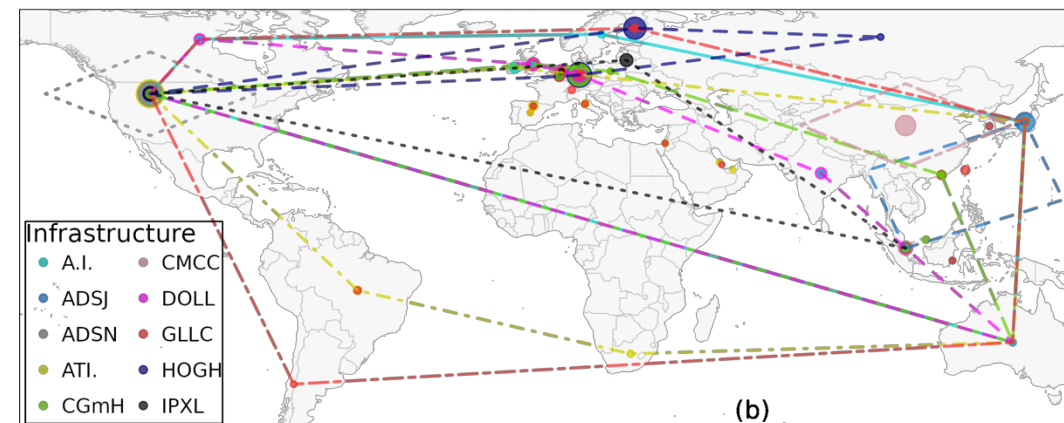
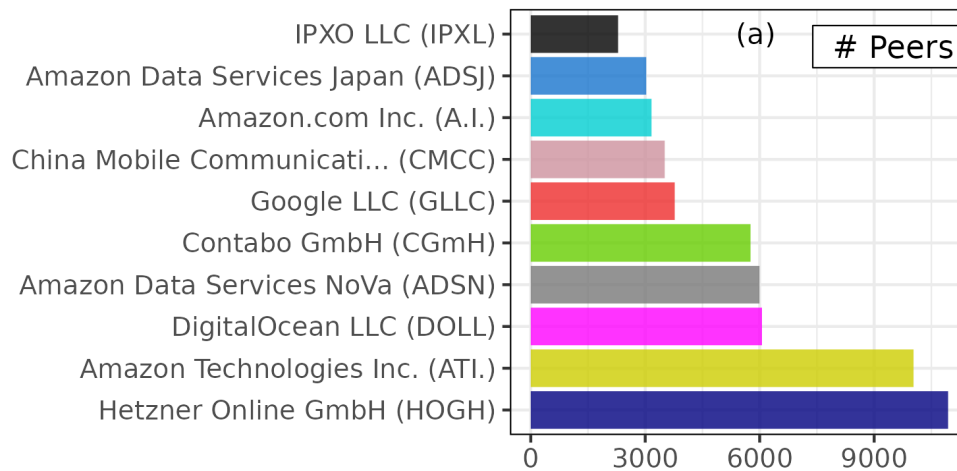
- Byzantine peers [Castro1999]
 - Consensus adversaries
 - Message adversaries
- The logs in replicas, i.e., ledgers, cope with **consensus adversaries**
 - The hashed-ordered chain of blocks avoids the double-spending [Nakamoto2008]
- Randomized iterative diffusion copes with oblivious **message adversaries**
 - Randomization and replenishing copes with churn [Maymounkov2002]
 - ~**System resilience**

Measurement Insights

The majority of the system relies on few well-known cloud providers

- **Infrastructure concentration** [Guzman2024-ICBC]

- **Low churn** (low frequency for peers going offline),
- High availability and reliability
- Internet-scale systems like IPFS, Bitcoin, TON, and XRP Ledger
- (Monetary) incentives motivate users to deploy in highly reliable and available infrastructures at the expense of **system resilience**



Measurement Insights

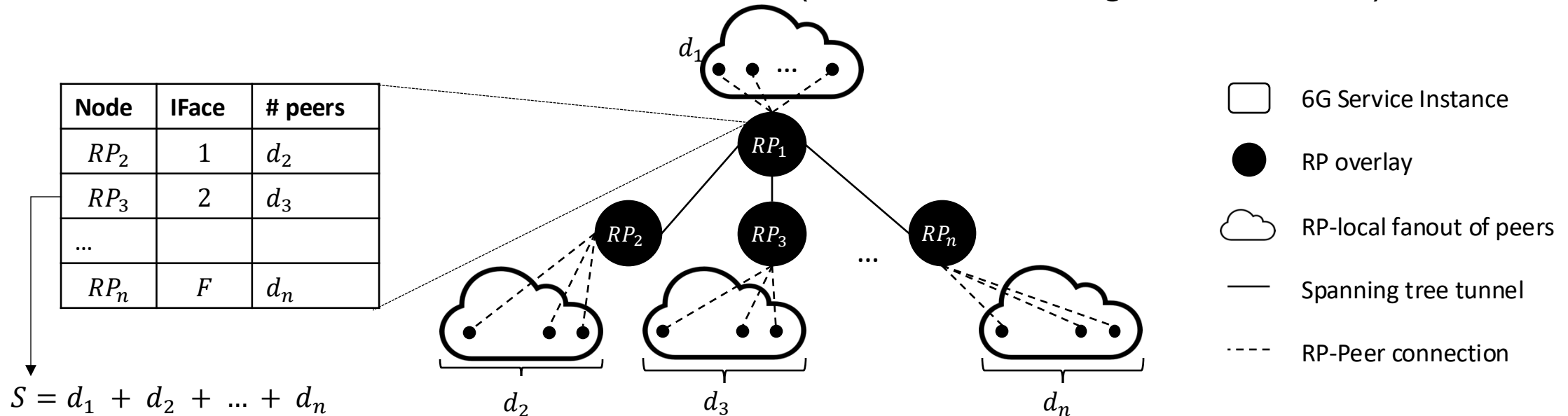
Internet relations from peer-to-peer to customer-provider

- The **delay** for opening a **communication relation follows a tail distribution** scale higher than the Internet delay, **~10x**, even though these distributed systems run in highly reliable infrastructures with few hops [Guzman2022-ICBC]
- Randomized iterative diffusion **cannot evaluate consensus finality**
 - Defining contention times to evaluate the consensus
 - Contention times are part of the proofs (e.g., PoX)
- **Expensive** communication costs for data and control planes due to the distributed nature of the consensus
 - DP: Distributed state, while randomly being flooded, gets **duplicated, 22 MB per 256 B transaction; on a day, 41 TB**
 - CP: **3 GB** is required only to discover and establish relations; this grows to a maintenance regime at a rate of **249 MB/s** [Guzman2024-DIN CoNEXT]

Is there Something better?

A multicasted approach

- Multicast-based diffusion
 - Places replication points (RPs) in infrastructure clusters
 - Replicates distributed state
 - RPs are not part of the consensus; consensus is still (randomized) **decentralized and fault tolerant** (oblivious message adversaries)



Is there Something better?

Multicast Insights

- There is no **delay** in opening a **communication relation**
 - **The diffusion can be enhanced by 5x** [Guzman2024-IFIP]
- Multicast **can evaluate consensus finality**
 - Informed contention times without expensive proofs (e.g., PoX)
- **Reduced** communication costs for data and control planes
 - DP: Distributed state, replicated with negligible extra cost
 - CP: Only **4 MB** is needed (**instead of 3 GB**) to establish relations with a maintenance regime 30x better than unicast [Guzman2024-DIN CoNEXT]